Initial Algebras Without Iteration

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Initial Algebras With Iteration

Initial Algebra Theorem (Trnková, Adámek, Koubek, Reiterman'75)

For a functor $F : Set \rightarrow Set$ tfae:

- 1. an initial algebra exists,
- 2. a fixed point exists,
- 3. a pre-fixed point exists,
- 4. the initial-algebra chain converges.

Initial-algebra chain: the unique ordinal-indexed chain

$$0 \to F0 \to FF0 \to \cdots \to W_{\omega} \to W_{\omega+1} \to \cdots \to W_{i} \xrightarrow{W_{i,i+1}} W_{i+1} \to \cdots$$
$$= \operatorname{colim} F^{i}0 = FW_{\omega} = FW_{i}$$

The chain converges if $w_{i,i+1}$ is an iso for some ordinal *i*.

more generally: category C with well-behaved class M of monos; F preserving M

$$(X \text{ with } FX \cong X)$$
$$(FX \xrightarrow{m} X \text{ monic})$$

Initial Algebras Without Iteration

Theorem

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 $(X \text{ with } FX \cong X)$ $(FX \xrightarrow{m} X \text{ monic})$

Our contribution: A new compact and constructive proof based on

- $\blacktriangleright\,$ Pataraia's fixed point theorem for dcpos with $\perp\,$
- pre-fixed points having smooth subobjects
- properties of recursive coalgebras

Ingredient 1: Pataraia's Theorem

Pataraia's and other Fixed Point Theorems

KleeneEvery continuous function on an ω -cpo with \perp has a least fixed point.= dcpo with \perp (Markowsky 1976)

Zermelo Every monotone function on a chain-complete poset has a least fixed point.

Knaster-Tarski Every monotone function on a complete lattice has least and greatest fixed point.

Theorem (Pataraia 1997)

Let *P* be a dcpo with \perp .

Then every monotone map $f: P \rightarrow P$ has a least fixed point μf .

Pataraia Induction Principle: Suppose $S \subseteq P$ fulfils

1.
$$\perp \in S$$
, 2. $s \in S \implies f(s) \in S$, 3. $\bigvee D \in S$ for $D \subseteq S$ directed.
Then $\mu f \in S$.

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Ingredient 2: Smooth of Subobjects and Monos

Smoothness

Definition

Let $\ensuremath{\mathcal{M}}$ be a class of monomorphisms closed under isos and composition.

- 1. The object A has smooth \mathcal{M} -subobjects if $\operatorname{Sub}_{\mathcal{M}}(A)$ is a dcpo with \bot where directed joins and \bot are given by colimits.
- 2. The class ${\mathcal M}$ is smooth if every object has smooth subobjects.

Directed joins and \perp given by colimits:

- ▶ initial object 0 exists, and $0 \rightarrow A$ lies in \mathcal{M} (empty colimit)
- ▶ for $D \subseteq \text{Sub}_{\mathcal{M}}(A)$ directed:



 $a_{u,v}$ witnesses $u \leq v$ in D $C = \operatorname{colim}_{u \in D} A_u$ $m = \bigvee D$ 1. (strong) monos in every lfp category with a simple initial object, e.g. sets, posets, graphs, monoids, vector spaces (but not rings)

all $0 \rightarrow A$ are strong monos

e embedding if $\exists \hat{e}. \ \hat{e} \cdot e = id$ and $e \cdot \hat{e} \leq id$.

- embeddings in dcpos with ⊥ + continuous maps (but all monos are not a smooth class)
- 3. (strong) monos in metric spaces + non-expansive maps

represent closed subspaces

 strong monos in complete metric spaces + non-expansive maps (but all monos are not a smooth class) Ingredient 3: Recursive Coalgebras

Recursive Coalgebras

Definition

A coalgebra $C \xrightarrow{\gamma} FC$ is recursive if for every algebra $FA \xrightarrow{\alpha} A$ there exists a unique coalgebra-to-algebra morphism

$$\begin{array}{ccc} C & \stackrel{h}{\longrightarrow} & A \\ \gamma \downarrow & & \uparrow^{\alpha} \\ FC & \stackrel{Fh}{\longrightarrow} & FA \end{array}$$

- ► Osius (1974): first studies them in connection with well-founded coalgebras for the power-set functor
- ► Taylor (1995, 1999, 2021): considers them for general functors as coalgebras obeying the recursion scheme; General Recursion Theorem: well-founded ⇒ recursive.
- Capretta, Uustalu and Vene (2006): constructions of recursive coalgebras; semantic treatment of recursive divide-and-conquer programs

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Key Properties of Recursive Coalgebras

- $0 \rightarrow F0$ is recursive
- $\blacktriangleright C \xrightarrow{\gamma} FC \text{ recursive} \implies FC \xrightarrow{F\gamma} FFC \text{ recursive}$
- recursive coalgebras are closed under all colimits (in the category of *F*-coalgebras)
- $\blacktriangleright \ C \xrightarrow{\gamma} FC \text{ recursive and } \gamma \text{ iso } \Longrightarrow FC \xrightarrow{\gamma^{-1}} C \text{ initial algebra}$

• Every coalgebra $W_i \xrightarrow{w_{i,i+1}} FW_i$ in the initial-algebra chain is recursive.

Combining the Ingredients

Theorem

Let $FA \xrightarrow{m} A$ be an \mathcal{M} -pre-fixed point of an endofunctor F preserving \mathcal{M} . If A has smooth \mathcal{M} -subobjects, then F has an initial algebra. (which is an \mathcal{M} -subalgebra of $FA \xrightarrow{m} A$)

Proof. Use Pataraia Induction on

$$f: \operatorname{Sub}_{\mathcal{M}}(A) \to \operatorname{Sub}_{\mathcal{M}}(A), \quad f(B \rightarrowtail^{u} A) = (FB \rightarrowtail^{Fu} FA \rightarrowtail^{m} A),$$

$$S = \left\{ B \xrightarrow{u} A : \exists B \xrightarrow{\beta} FB \text{ recursive s.th.} \begin{array}{c} B \xrightarrow{u} A \\ \beta \downarrow & \uparrow m \\ FB \xrightarrow{Fu} FA \end{array} \right\}.$$

all $u \in \operatorname{Sub}_{\mathcal{M}}(A) \text{ s.th. } u \leq f(u)$
via a recursive coalgebra

Initial Algebra Theorem (Proof continued)

$$S = \{B \xrightarrow{u} A : \exists B \xrightarrow{\beta} FB \text{ recursive s.th. } u = f(u) \cdot \beta.\}$$

- 1. $\bot \in S$: clearly $u: 0 \rightarrow A$ is in S
- 2. S closed under f: if $u \in S$, then $f(u) \in S$ since

 $f(u) = m \cdot Fu = m \cdot F(f(u) \cdot \beta) = m \cdot F(f(u)) \cdot F\beta = f(f(u)) \cdot F\beta$ (\$\beta\$ recursive \Rightarrow F\beta\$ recursive)

3. S closed under directed joins: given $D \subseteq S$ directed, $u \leq v$ in D is necessarily witnessed by a coalgebra homomorphism $h: B_u \rightarrow B_v$:



Now let $v = \bigvee D$. Then

- $\blacktriangleright B_v = \operatorname{colim}_{u \in D} B_u$
- v uniquely induced by colimit
- $\blacktriangleright B_{\nu} \xrightarrow{\beta_{\nu}} FB_{\nu} \text{ recursive}$
- ► v coalgebra-to-algebra morphism

Therefore $v \in S$.

Initial Algebra Theorem (Proof continued)

- ▶ By Parataia Induction: $\mu f \in S$, ... let's denote it $I \xrightarrow{u} A$.
- There is a recursive coalgebra $I \xrightarrow{\iota} FI$ witnessing $u \leq f(u)$:



- Since u = f(u) in $Sub_{\mathcal{M}}(A)$, ι is an iso.
- Therefore $FI \xrightarrow{\iota^{-1}} I$ is an initial algebra.

Initial Algebra Theorem

Corollary

Let \mathcal{C} have a smooth class \mathcal{M} of monomorphisms. For $F: \mathcal{C} \to \mathcal{C}$ preserving \mathcal{M} tfae:

- ▶ an initial algebra exists,
- ▶ a fixed point exists,
- \blacktriangleright an $\mathcal M\text{-}\mathsf{pre}\text{-}\mathsf{fixed}$ point exists.

Moreover, if these hold, then μF is an \mathcal{M} -subalgebra of every \mathcal{M} -pre-fixed point of F.

 A constructive proof of Trnková et al.'s initial algebra theorem. (omitting the transfinite initial-algebra chain)

In the paper:

- A streamlined (non-constructive) proof of the original theorem (with initial-algebra chain) based on Zermelo's theorem.
- ► An application to categories enriched in dcpos with ⊥: coincidence of initial algebra and terminal coalgebra, if they exist.

Appendix ...

Dcpo Enriched Categories

Corollary

Let C be enriched in the category of dcpos with \perp and continuous maps with composition strict in both arguments. For $F: C \rightarrow C$ locally monotone tfae:

- ▶ an initial algebra exists,
- ▶ a terminal coalgebra exists,
- ▶ a fixed point exists,
- ▶ a pre-fixed point carried by an embedding exist.

Moreover: $FI \xrightarrow{\iota} I$ initial algebra $\implies I \xrightarrow{\iota^{-1}} FI$ terminal coalgebra.

Theorem

Let C be enriched in the category of dcpos with \perp and continuous maps with composition left-strict. Let $F: C \to C$ be locally monotone.

Then: $FI \xrightarrow{\iota} I$ initial algebra $\implies I \xrightarrow{\iota^{-1}} FI$ terminal coalgebra.

Freyd proved this for locally continuous functors using Kleene's theorem. Here we use Pataraia Induction.